## Flexible techniques for analyzing and manipulating web service protocols

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Material based on work in **ServiceMosaic** project (http://servicemosaic.isima.fr )

Collaboration with colleagues and students

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- Dr. Fabio Casati (HP Labs, USA)
- Mr. Hamid Motahari (UNSW, Australia)
- Mr. Julien Ponge (LIMOS, France)

#### Agenda

- Web services vision and technologies
- Representing web service protocols
- Analysis and management of web service protocols
- Summary and outlook

#### Motivations

- Integration of autonomous and heterogeneous systems
- Automation of inter-organizational business processes
- Web services: evolution of current technologies
- Distributed information systems
- Middleware (RPC, MOM, CORBA, ),
- Entreprise Application Integration

## Beyond current technologies

- New integration context
- Open environment : autonomous systems
- Large and dynamic integration space
- Semantic heterogeneity (both data and business processes): one-to-one mappings between partner systems do not scale
- Inter-organizational interactions (trust, security, transactions,
- Limitations of current technologies
- Centralized middleware
- Rigid infrastructures, costly development and maintenance of integrated systems
- Close environment/tightly coupled systems (semantics known from the context)

#### Web services

protocols" [W3C] agents using XML-based messages exchanged via Internet-based bindings are capable of being defined, described and discovered as XML artifacts. A web service supports direct interactions with other software "a software application identified by a URI, whose interfaces and

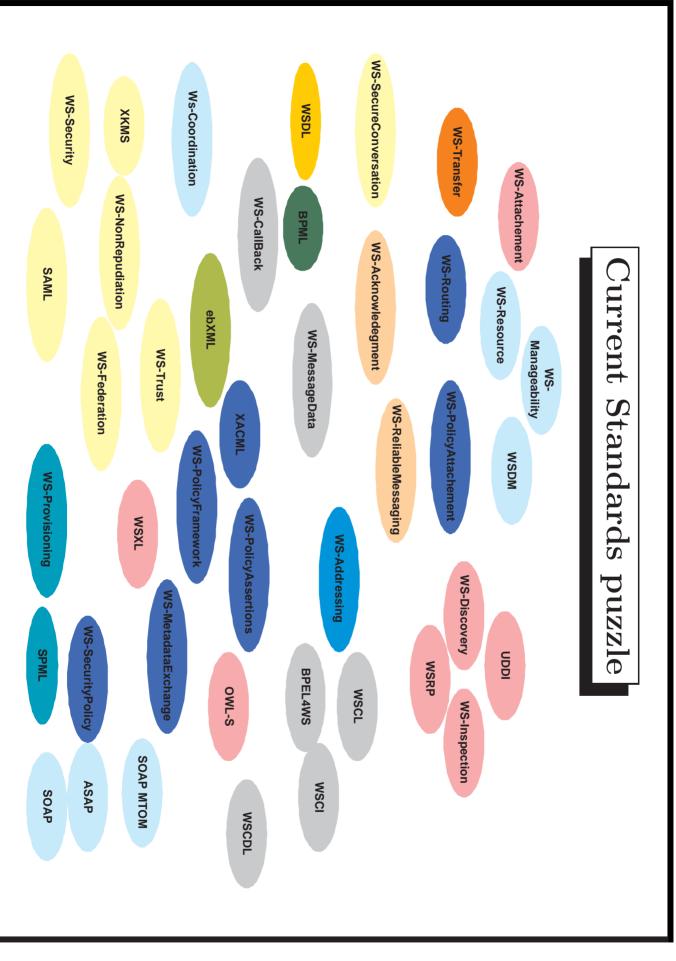
### Main characteristics

- Generic interface for service oriented architectures
- Intensive use of standards (SOAP, WSDL, etc)
- Loosely coupled integration
- of distributed applications ⇒ Ultimate goal: rapid low-cost development and easy composition

## Web service technologies

Specifications and languages providing core functionality of web services

- Service description
- Service discovery
- Service interactions
- Service composition



### Interoperability layers

Policy Layer

Business Interface &

Protocol

Layer

functional properties (e.g., cost, response time, ..) Policy specification (e.g., privacy policies) and non-

Basic Coordination

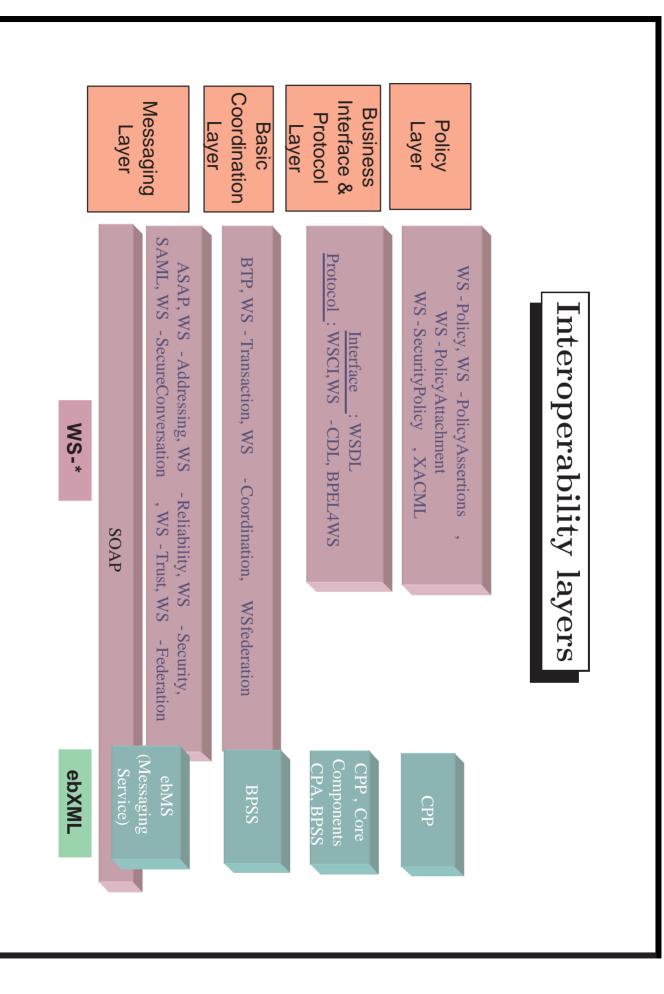
Functional properties of services (interfaces and business protocols)

Requirements and properties related to a set of

message exchanges among two or more partners

Messaging Layer

Standard information transportation protocol



### Some Observations

- Services are loosely-coupled and need to be fully specified explicit (essential in autonomous environments) Making implicit information (as in closed environments)
- Interface,
- Business protocols,
- Functional and non-functional properties (e.g., QoS, ...),
- Meaning of the parameters, operations effects, negociation parameters, ...
- Trade-off: expressive power vs. readability/usability

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### ServiceMosaic project

Joint project with SOC group (UNSW, Sydney) and HP laboratories (Palo Alto, USA)

- Definition of a service description framework
- semantics A protocol model endowed with richer abstractions and a formal
- Design of an algebra for high level analysis and management of service protocols
- Protocol discovery
- Model-driven approach to support service adaptation
- Model-driven change impacts analysis
- Development of a fully-fledged CASE tool for Web service development and lifecycle management

### Protocol modeling

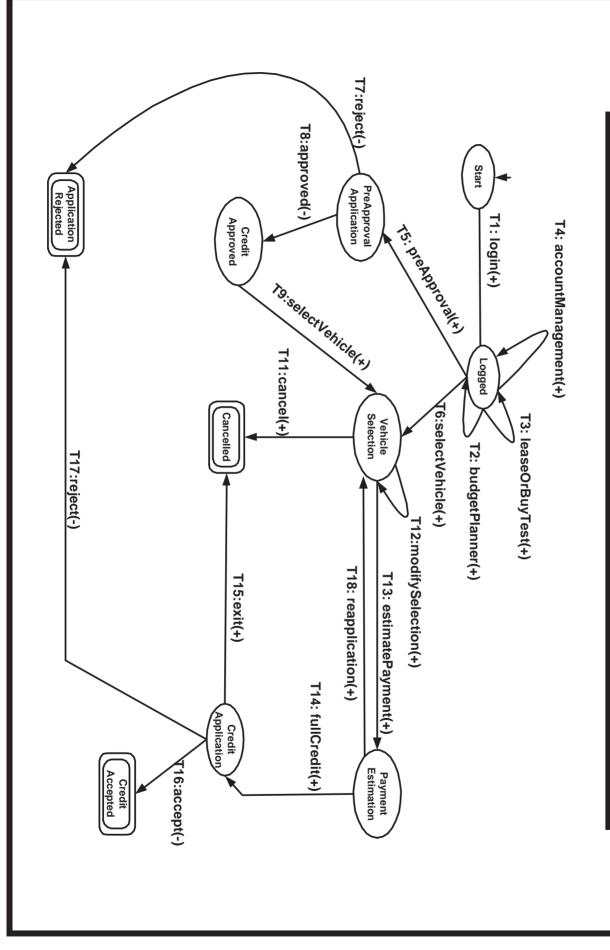
Describe external behavior of services

An extended protocol model

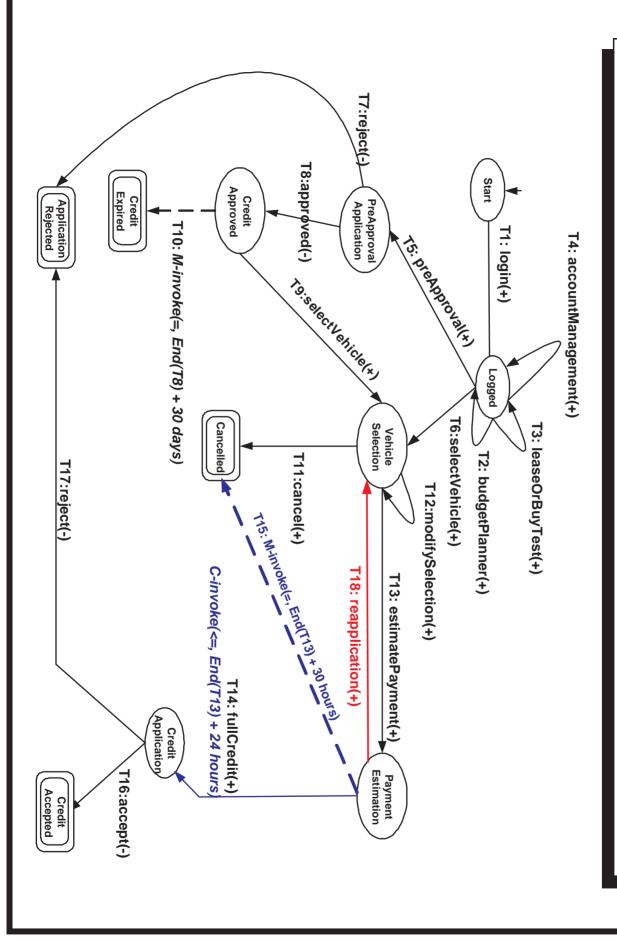
- Message Choreography
- Time-sensitive Conversations
- Transactional Implications and Effects

Rightsizing is a key issue

# Example: a credit online financing services



# Example: a credit online financing services (cont.)



# Protocol analysis and management

- Two dimensions of the analysis
- Compatibility: checking whether two services can interact correctly based on their protocol definitions
- Replaceability: verifying whether two protocols can support the same set of conversations
- Characterization of different levels of protocol compatibility and replaceability
- Need for a protocol algebra
- primitives to analyze and manage protocols
- key to achieve the benefits
- tools that implement these operators
- and replaceability between pairs of protocols ⇒ Provide an automated support for verification of compatibility

## Applications and benefits

- Service discovery and composition (e.g., reduce the number of false positives during service discovery)
- Change support and evolution
- Support for static and dynamic binding
- Compliance verification (e.g., with B2B standards)

#### Formalization

### !! Small is beautiful

A basic protocol: message choreography + message polarity

- Protocol model based on a finite state machine
- Protocol semantics: branching time view

behavior) Protocol  $\mathcal{P}$  = a schema (i.e., intentional description of service

Semantics of a protocol: set of complete execution trees

Semantics of interactions: set of complete interaction trees

- $\Rightarrow$  Useful to deal with service composition
- A slightly adapted simulation relation to compare protocols

### Compatibility classes

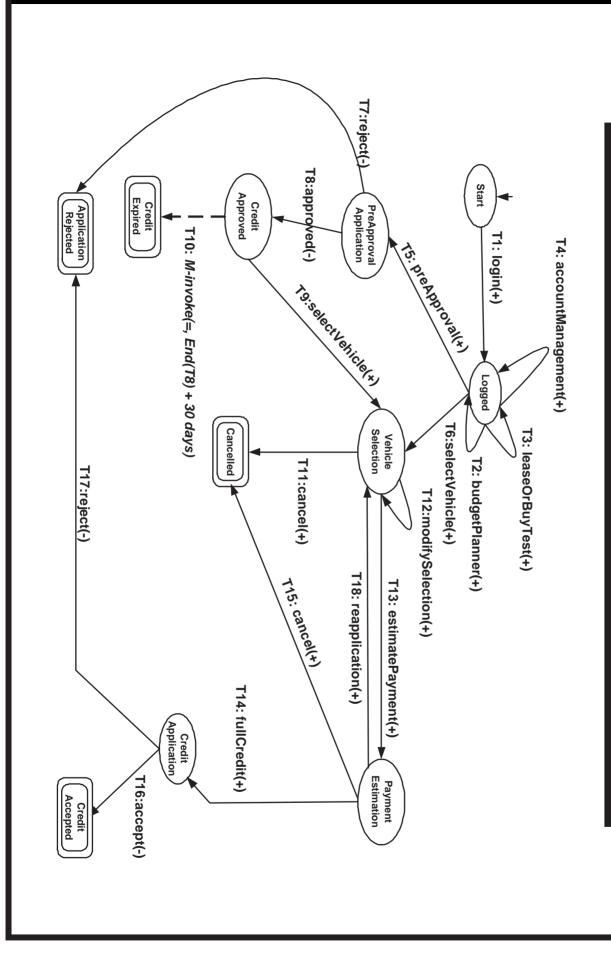
### Partial compatibility

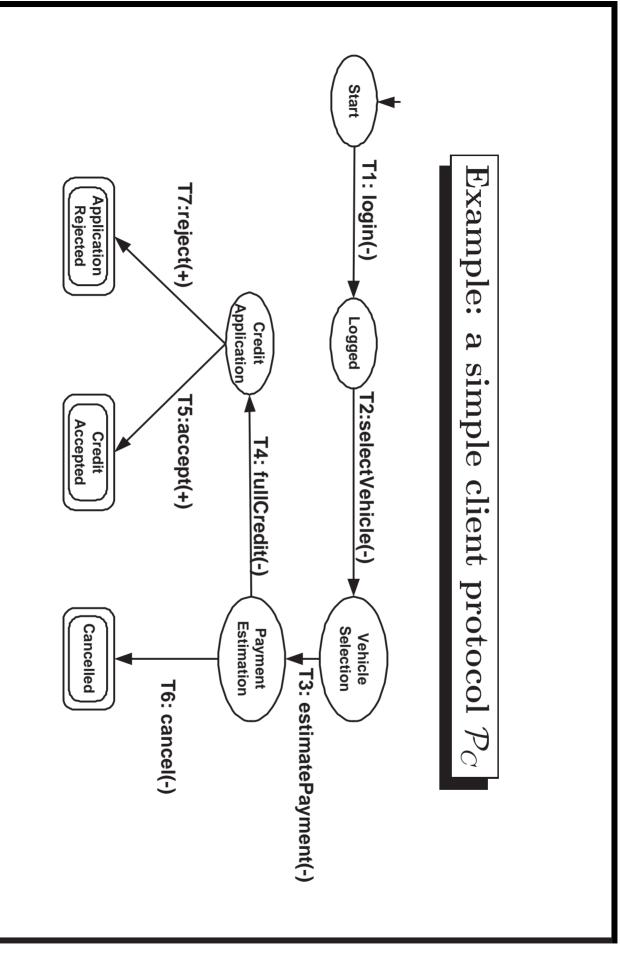
possible conversation that can take place  $\mathcal{P}_1$  that can interoperate with  $\mathcal{P}_2$ , i.e., if there is at least one  $\mathcal{P}_1$  is partially compatible with  $\mathcal{P}_2$  if there are some executions of

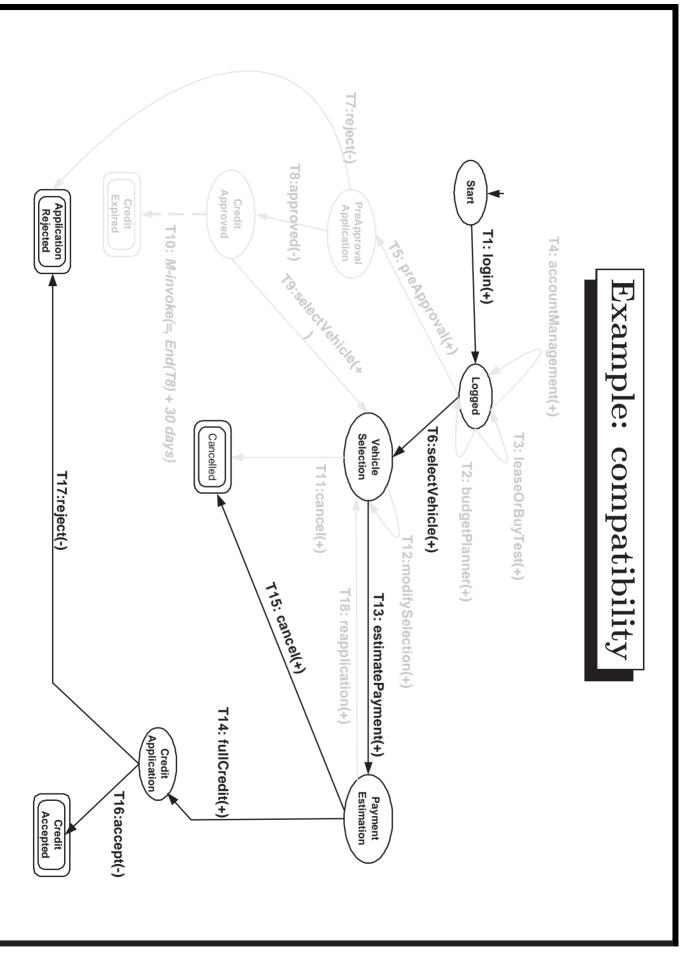
#### Full compatibility

 $\mathcal{P}_1$  is fully compatible with  $\mathcal{P}_2$  if all executions of  $\mathcal{P}_1$  can by  $\mathcal{P}_1$  is understood by  $\mathcal{P}_2$ interoperate with  $\mathcal{P}_2$ , i.e., any conversation that can be generated

# Example: a basic online financing services



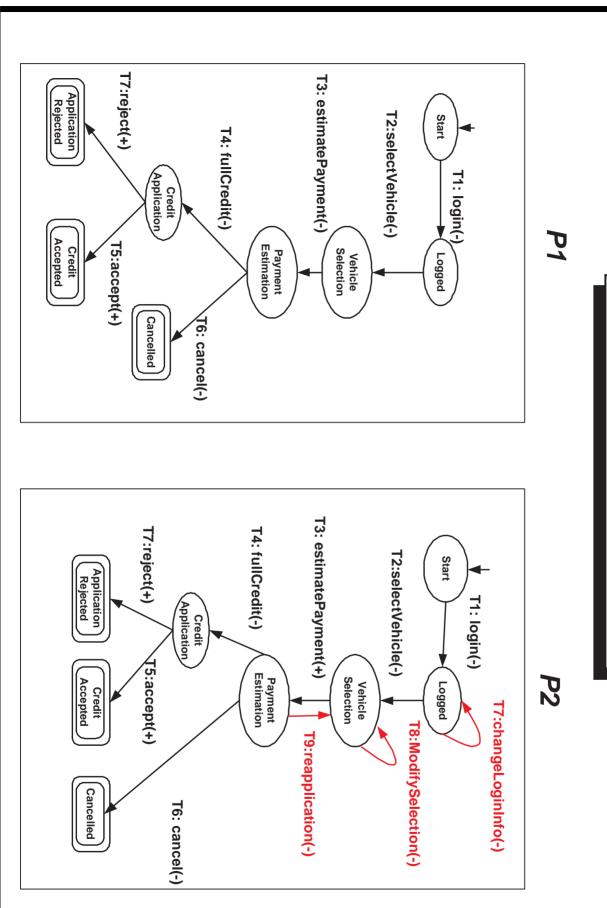




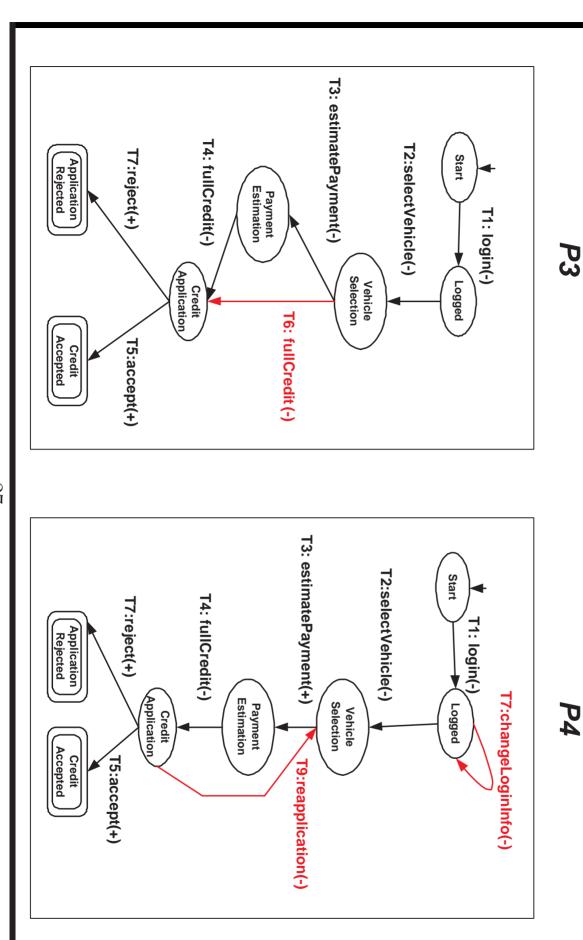
### Replaceability classes

- Protocol equivalence
- context and the change is transparent to clients Identifies when two protocols can be interchangeably used in any
- Protocol subsumption
- $\mathcal{P}_2$  (the opposite is not necessarily true) Identifies when a protocol  $\mathcal{P}_1$  can be transparently used instead of
- Protocol equivalence and subsumption with respect to a client protocol
- Identifies replaceability relations with respect to a certain client
- Protocol equivalence and subsumption with respect to an interaction role
- $\mathcal{P}_2$  can replace  $\mathcal{P}_1$  with respect to a role  $\mathcal{P}_R$  if  $\mathcal{P}_2$  behaves as  $\mathcal{P}_1$ when  $\mathcal{P}_1$  behaves as  $\mathcal{P}_R$

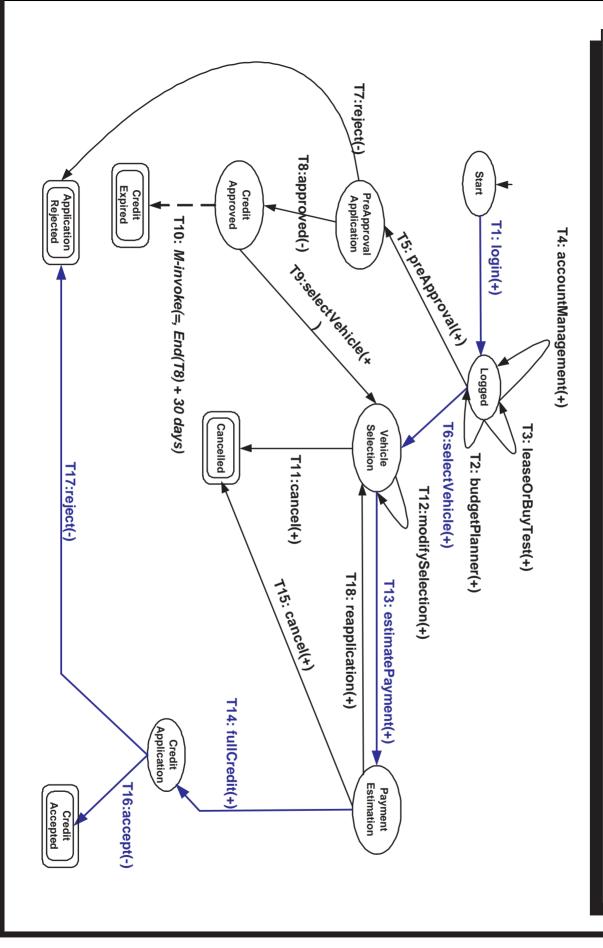
## Example: subsumption



### Example: equivalence w.r.t. a client protocol



## Example: equivalence w.r.t. a client protocol (cont.)



## Toward a protocol algebra

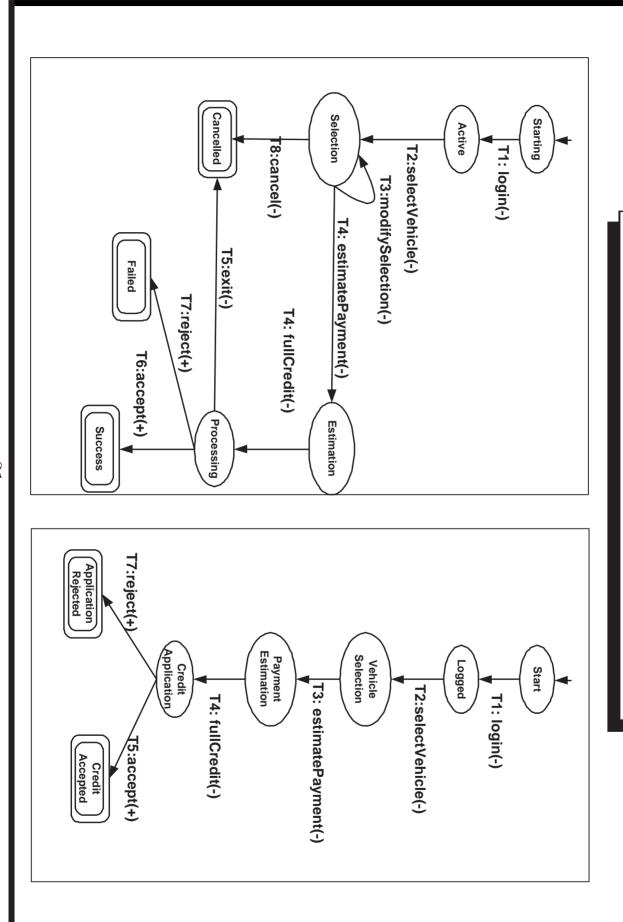
A set of operators to manipulate and analyze protocols

#### Example

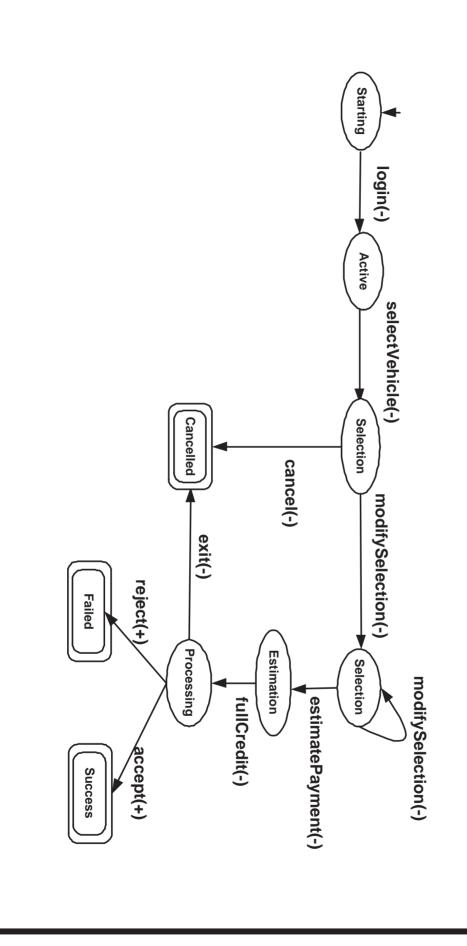
- Compatible composition
- Returns a protocol that describes all the possible conversations between the considered protocols
- Intersection
- common between the input protocols Returns a protocol that describes the set of conversations that are
- Difference
- protocol first input protocol that cannot be supported by the second input Returns a protocol that describes the set of conversations of the

#### T4: fullCredit(-) estimateP T2:selectVehicle(-) ayment(-) T7:reject(+) T1: login(-) Example: compatible composition fullCredit (-) T5:accept(+) **T**6: estimatePayment(+) T2:selectVehicle(+) T4: fullCredit(+) T7:reject(-) T1: login(+) T5:accept(-) T9:reappli cation(+) reject login selectVehicle **fullCredit** accept **Payment** estimate

## Example: difference operation



# Example: difference operation (cont.)



## Characterization of the classes

Partial compatibility

P-compat $(\mathcal{P}_1, \mathcal{P}_2)$  iff  $\mathcal{P}_1||^{c}\mathcal{P}_2$  is not an empty protocol

• Full compatibility

F-compat $(\mathcal{P}_1, \mathcal{P}_2)$  iff  $[\mathcal{P}_1||^{\mathsf{c}}\mathcal{P}_2]_{\mathcal{P}_1} \cong \mathcal{P}_1$ .

Subsumption

 $Subs(\mathcal{P}_2, \mathcal{P}_1) \text{ iff } \mathcal{P}_2 \lesssim \mathcal{P}_1$ 

Equivalence

 $Equiv(\mathcal{P}_1, \mathcal{P}_2) \text{ iff } \mathcal{P}_1 \cong \mathcal{P}_2$ 

Replaceability with respect to a client protocol  $Repl_{[\mathcal{P}_c]}(\mathcal{P}_1, \mathcal{P}_2)$  iff  $\mathcal{P}_c||^{\mathsf{c}}(\mathcal{P}_2||^{\mathsf{d}}\mathcal{P}_1)$  is an empty protocol

Replaceability with respect to an interaction role  $Repl\_Role_{[\mathcal{P}_R]}(\mathcal{P}_1, \mathcal{P}_2) \text{ iff } (P_R||^{\mathsf{I}}\mathcal{P}_2) \lesssim \mathcal{P}_1.$ 

## Extension to timed protocols

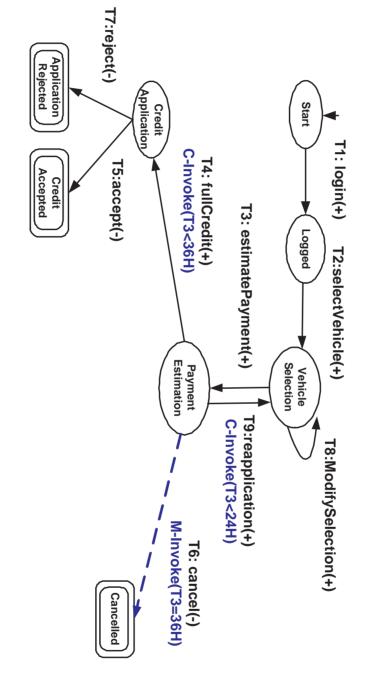
Timed protocol = basic protocol + temporal abstractions

- Two main temporal abstractions
- C-invoke: temporal window
- M-Invoke: expiration
- Formal semantics

Protocol semantics: set of timed execution paths (conversation)

Interaction semantics: set of timed interaction paths

## Example of a timed protocol

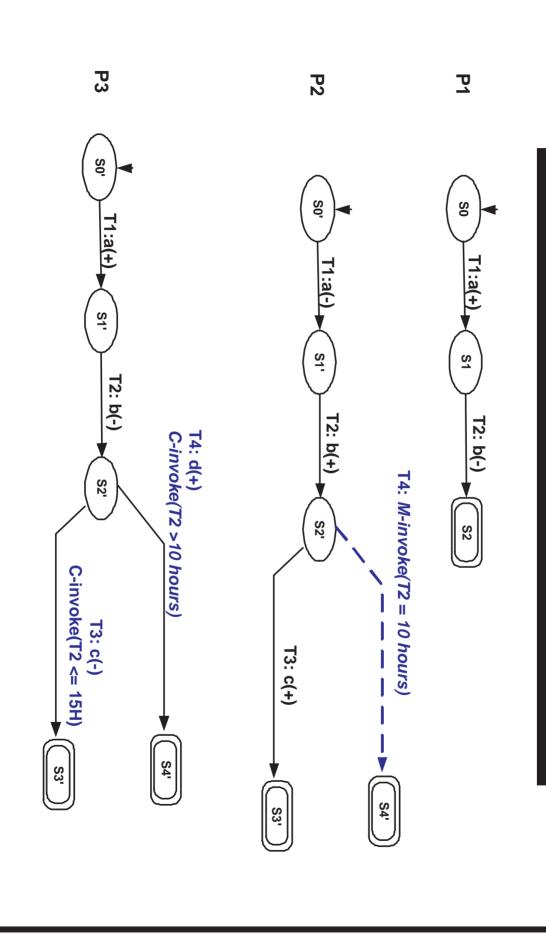


Timed conversation: (login(+), 0); (select Vehicle(+), 1);(estimatePayment(+), 2); (cancel(-), 40)

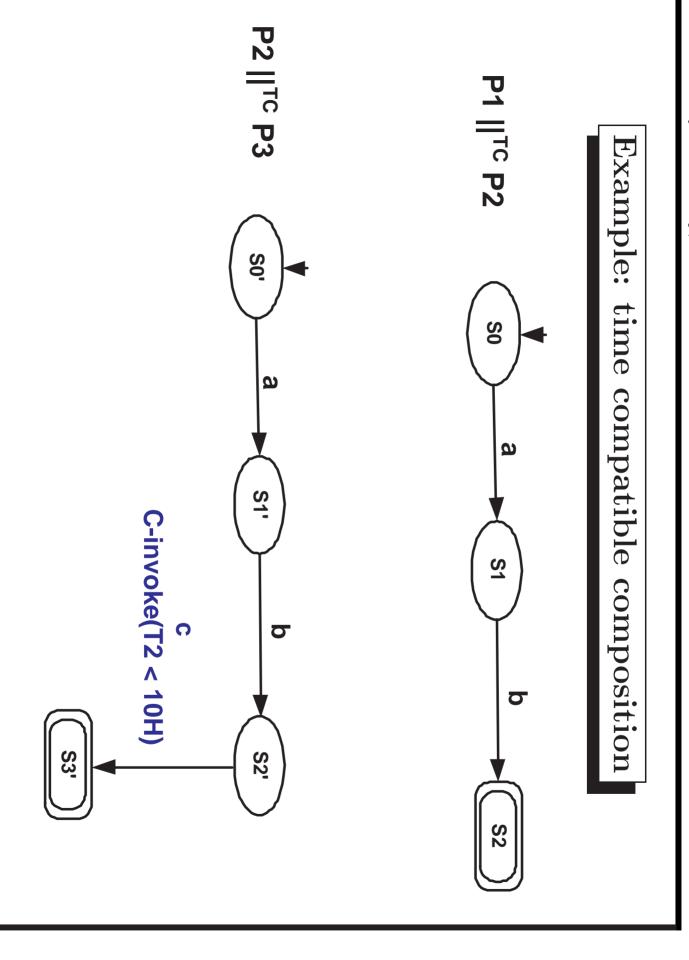
# Analysis and management of timed protocols

- Definition of new time-sensitive replaceability and compatibility
- Definition of operators for analyzing and managing timed protocols

# Example: time-sensitive compatibility



F.Toumani- QSL - Nancy, 27 June 2006

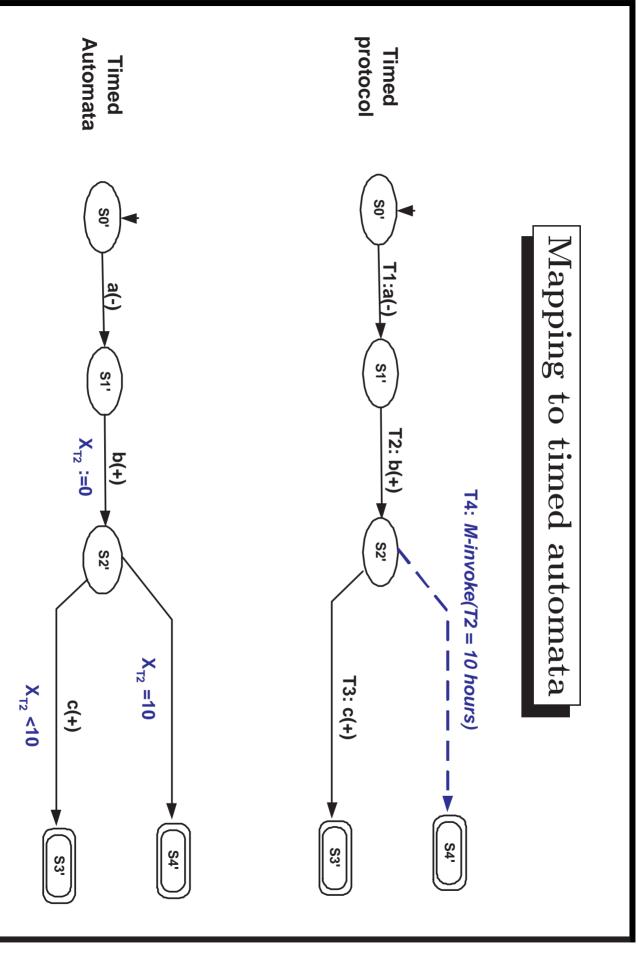


# Timed protocols: decision problems

# Strong link with the theory of timed automata

- Formal notation to model behavior of real time systems
- State-transition graphs with timing contraints using real-valued clock variables
- Extensively studied formalism

Mapping: timed protocols  $\rightarrow$  timed automata



#### Main results

A few lessons from Timed Automata

- Closed under all boolean operations but complementation
- clocks Silent transitions strictly increase expressiveness when they reset

Timed protocols = New subclass of timed automata

 $\rightarrow$  closed under complementation

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## Summary and outlook

- Analysis and management of web service protocols development and interoperability So far, focus on primitives for facilitating automation of services
- Extension to business protocols augmented with transactional abstractions
- Trust negotiation and security protocols in Web services
- Analysis of multi-party protocols (consistency analysis)
- Composite services analysis and synthesis
- Protocol discovery