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Time-stamp: < 10 Jun 2002 at 14:06:57 by charpov on berlioz.cs.unh.edu >
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Non blocking Atomic Committment Protocol (ACP-NB)

The non blocking property AC5 is obtained by using a reliable broadcast implemented as follows:

- upon reception of a broadcast message, this message is forwarded to all participants before it's delivered to the local site;
- since participant i does not forward to itself, $\mathit{forward}[i]$ is used to
- store the decision before it's delivered (and becomes "decision")

EXTENDS ACP_SB

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Participants type is extended with a "forward" variable.

Coordinator type is unchanged.

TypeInvParticipantNB \triangleq participant \in [

    participants \rightarrow [

    vote : {yes, no},

    alive : BOOLEAN,

    decision : {undecided, commit, abort},

    faulty : BOOLEAN,

    voteSent : BOOLEAN,

    forward : [participants \rightarrow {notsent, commit, abort}]

]
```

 $TypeInvNB \triangleq TypeInvParticipantNB \land TypeInvCoordinator$

Initially, participants have not forwarded anything yet

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InitParticipantNB \triangleq participant \in [
participants \rightarrow [
vote : {yes, no},
alive : {TRUE},
decision : {undecided},
faulty : {FALSE},
voteSent : {FALSE},
forward : [participants \rightarrow {notsent}]
]
```

 $InitNB \triangleq InitParticipantNB \land InitCoordinator$

Participant statements that realize a better broadcast

forward(i, j): forwarding of the predecision from participant i to participant jIF participant i is alive participant i has received a decision (stored in forward[i]) participant i has not yet forwarded this decision to participant jTHEN

participant i forwards the decision to participant j

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 \begin{bmatrix} @ \text{ EXCEPT } !.forward = \\ & [@ \text{ EXCEPT } ![j] = participant[i].forward[i]] \\ & ] \\ & ] \\ & \land \text{ UNCHANGED } \langle coordinator \rangle
```

preDecideOnForward(i, j): participant i receives decision from participant j

IF participant i is alive participant i has yet to receive a decision participant j has forwarded its decision to participant iTHEN

participant i (pre)decides in accordance with participant j's decision

 $\begin{aligned} preDecideOnForward(i, j) &\triangleq \land i \neq j \\ \land participant[i].alive \\ \land participant[i].forward[i] = notsent \\ \land participant[j].forward[i] \neq notsent \\ \land participant' = [participant EXCEPT ![i] = \\ & [@ EXCEPT !.forward = \\ & [@ EXCEPT ![i] = participant[j].forward[i]] \\ &] \\ &] \\ \land \text{UNCHANGED } \langle coordinator \rangle \end{aligned}$

preDecide(i): participant *i* receives decision from coordinator

IF participant i is alive participant i has yet to receive a decision coordinator has sent its decision to participant iTHEN participant i (pre)decides in accordance with coordinator's decision

 $preDecide(i) \triangleq \land participant[i].alive$

decideNB(i): Actual decision, after predecision has been forwarded IF participant *i* is alive participant *i* has forwarded its (pre)decision to all other participants

THEN

participant i decides in accordance with it's predecision

 $\begin{array}{ll} decideNB(i) \triangleq & \land participant[i].alive \\ & \land \forall j \in participants : participant[i].forward[j] \neq notsent \\ & \land participant' = [participant \ \texttt{EXCEPT} \ ![i] = \\ & & [@ \ \texttt{EXCEPT} \ !.decision = participant[i].forward[i]] \\ & & \\ & & \\ & \land \texttt{UNCHANGED} \ \langle coordinator \rangle \end{array}$

abortOnTimeout(i): conditions for a timeout are simulated IF

participant is alive and undecided and coordinator is not alive

coordinator died before sending decision to all participants who are alive all dead participants died before forwarding decision to a participant who is alive THEN decide abort

 $abortOnTimeout(i) \triangleq \land participant[i].alive \\ \land participant[i].decision = undecided \\ \land \neg coordinator.alive \\ \land \forall j \in participants : participant[j].alive \Rightarrow coordinator.broadcast[j] = notsent \\ \land \forall j, k \in participants : \neg participant[j].alive \land participant[k].alive \Rightarrow participant[j].forward[k] = notsent \\ \land participant' = [participant EXCEPT ![i] = [@ EXCEPT !.decision = abort]] \\ \land UNCHANGED \langle coordinator \rangle$

FOR N PARTICIPANTS

 $parProgNB(i, j) \triangleq \lor sendVote(i) \\ \lor abortOnVote(i) \\ \lor abortOnTimeoutRequest(i) \\ \lor forward(i, j) \\ \lor preDecideOnForward(i, j) \\ \lor abortOnTimeout(i) \\ \lor preDecide(i) \\ \lor decideNB(i)$

 $parProgNNB \triangleq \exists i, j \in participants : parDie(i) \lor parProgNB(i, j)$

 $progNNB \triangleq parProgNNB \lor coordProgN$

 $\begin{array}{l} fairnessNB \ \triangleq \ \land \forall i \in participants : WF_{\langle coordinator, participant \rangle}(\exists j \in participants : parProgNB(i, j)) \\ \land WF_{\langle coordinator, participant \rangle}(coordProgB) \end{array}$

 $SpecNB \stackrel{\Delta}{=} InitNB \land \Box [progNNB]_{(coordinator, participant)} \land fairnessNB$

(SOME) INVALID PROPERTIES

 $\begin{aligned} AllCommit &\triangleq \forall i \in participants : \diamondsuit(participant[i].decision = commit \lor participant[i].faulty) \\ AllAbort &\triangleq \forall i \in participants : \diamondsuit(participant[i].decision = abort \lor participant[i].faulty) \\ AllCommitYesVotes &\triangleq \forall i \in participants : & \forall j \in participants : participant[j].vote = yes \\ &\sim participant[i].decision = commit \lor participant[i].faulty \lor coordinator.faulty \end{aligned}$